

15.093 Optimization Methods

Lecture 20: The Conjugate Gradient Algorithm
Optimality conditions for constrained optimization

1 Outline

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1. The Conjugate Gradient Algorithm
2. Necessary Optimality Conditions
3. Sufficient Optimality Conditions
4. Convex Optimization
5. Applications

2 The Conjugate Gradient Algorithm

2.1 Quadratic functions

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$$\min f(\mathbf{x}) = \frac{1}{2} \mathbf{x}' \mathbf{Q} \mathbf{x} + \mathbf{c}' \mathbf{x}$$

Definition: $\mathbf{d}_1, \dots, \mathbf{d}_n$ are \mathbf{Q} -conjugate if

$$\mathbf{d}_i \neq \mathbf{0}, \quad \mathbf{d}_i' \mathbf{Q} \mathbf{d}_j = 0, \quad i \neq j$$

Proposition: If $\mathbf{d}_1, \dots, \mathbf{d}_n$ are \mathbf{Q} -conjugate, then $\mathbf{d}_1, \dots, \mathbf{d}_n$ are linearly independent.

2.2 Motivation

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Given \mathbf{Q} -conjugate $\mathbf{d}_1, \dots, \mathbf{d}_n$, and \mathbf{x}^k , compute

$$\begin{aligned} \min_{\alpha} f(\mathbf{x}^k + \alpha \mathbf{d}_k) &= \mathbf{c}' \mathbf{x}^k + \alpha \mathbf{c}' \mathbf{d}_k + \\ &\frac{1}{2} (\mathbf{x}^k + \alpha \mathbf{d}_k)' \mathbf{Q} (\mathbf{x}^k + \alpha \mathbf{d}_k) = \\ &f(\mathbf{x}^k) + \alpha \nabla f(\mathbf{x}^k)' \mathbf{d}_k + \frac{1}{2} \alpha^2 \mathbf{d}_k' \mathbf{Q} \mathbf{d}_k \end{aligned}$$

Solution:

$$\hat{\alpha}_k = \frac{-\nabla f(\mathbf{x}^k)' \mathbf{d}_k}{\mathbf{d}_k' \mathbf{Q} \mathbf{d}_k}, \quad \mathbf{x}^{k+1} = \mathbf{x}^k + \hat{\alpha}_k \mathbf{d}_k$$

2.3 Expanding Subspace Theorem

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$\mathbf{d}_1, \dots, \mathbf{d}_n$ are \mathbf{Q} -conjugate. Then, \mathbf{x}^{k+1} solves

$$\begin{aligned} \min \quad &f(\mathbf{x}) \\ \text{s.t.} \quad &\mathbf{x} = \mathbf{x}_1 + \sum_{j=1}^k \alpha_j \mathbf{d}_j \end{aligned}$$

Moreover, $\mathbf{x}^{n+1} = \mathbf{x}^*$.

2.4 The Algorithm

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Step 0 Given \mathbf{x}^1 , set $k := 1$, $\mathbf{d}_1 = -\nabla f(\mathbf{x}^0)$

Step 1 For $k = 1, \dots, n$ do:

If $\|\nabla f(\mathbf{x}^k)\| \leq \epsilon$, stop; else:

$$\hat{\alpha}_k = \operatorname{argmin}_{\alpha} f(\mathbf{x}^k + \alpha \mathbf{d}^k) = \frac{-\nabla f(\mathbf{x}^k)' \mathbf{d}_k}{\mathbf{d}'_k Q \mathbf{d}_k}$$

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \hat{\alpha}_k \mathbf{d}_k$$

$$\mathbf{d}_{k+1} = -\nabla f(\mathbf{x}^{k+1}) + \lambda_k \mathbf{d}_k, \quad \lambda_k = \frac{-\nabla f(\mathbf{x}^{k+1})' Q \mathbf{d}_k}{\mathbf{d}'_k Q \mathbf{d}_k}$$

2.5 Correctness

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Theorem: The directions $\mathbf{d}_1, \dots, \mathbf{d}_n$ are Q -conjugate.

2.6 Convergence Properties

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- This is a finite algorithm.
- If there are only k distinct eigenvalues of Q , the CGA finds an optimal solution in k steps.
- Idea of pre-conditioning. Consider

$$\min f(S\mathbf{x}) = \frac{1}{2}(S\mathbf{x})' Q(S\mathbf{x}) + \mathbf{c}' S\mathbf{x}$$

so that the number of distinct eigenvalues of $S'QS$ is small

2.7 Example

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$$f(\mathbf{x}) = \frac{1}{2}\mathbf{x}' Q \mathbf{x} - \mathbf{c}' \mathbf{x}$$

$$Q = \begin{pmatrix} 35 & 19 & 22 & 28 & 16 & 3 & 16 & 6 & 4 & 4 \\ 19 & 43 & 33 & 19 & 5 & 2 & 5 & 4 & 0 & 0 \\ 22 & 33 & 40 & 29 & 12 & 7 & 6 & 2 & 2 & 4 \\ 28 & 19 & 29 & 39 & 16 & 7 & 14 & 6 & 2 & 4 \\ 16 & 5 & 12 & 16 & 12 & 4 & 8 & 2 & 4 & 8 \\ 3 & 2 & 7 & 7 & 4 & 5 & 1 & 0 & 1 & 4 \\ 16 & 5 & 6 & 14 & 8 & 1 & 12 & 2 & 2 & 4 \\ 6 & 4 & 2 & 6 & 2 & 0 & 2 & 4 & 0 & 0 \\ 4 & 0 & 2 & 2 & 4 & 1 & 2 & 0 & 2 & 4 \\ 4 & 0 & 4 & 4 & 8 & 4 & 4 & 0 & 4 & 16 \end{pmatrix} \quad \text{and} \quad \mathbf{c} = \begin{pmatrix} -1 \\ 0 \\ 0 \\ -3 \\ 0 \\ -2 \\ 0 \\ -6 \\ -7 \\ -4 \end{pmatrix}$$

$$\kappa(Q) \approx 17,641 \quad \delta = \left(\frac{\kappa(Q)-1}{\kappa(Q)+1} \right)^2 = 0.999774 \quad \text{SLIDE 9}$$

k	$f(\mathbf{x}^k)$	$f(\mathbf{x}^k) - f(\mathbf{x}^*)$	$\frac{f(\mathbf{x}^k) - f(\mathbf{x}^*)}{f(\mathbf{x}^{k-1}) - f(\mathbf{x}^*)}$
1	12.000000	2593.726852	1.000000
2	8.758578	2590.485430	0.998750
3	1.869218	2583.596069	0.997341
4	-12.777374	2568.949478	0.994331
5	-30.479483	2551.247369	0.993109
6	-187.804367	2393.922485	0.938334
7	-309.836907	2271.889945	0.949024
8	-408.590428	2173.136424	0.956532
9	-754.887518	1826.839334	0.840646
10	-2567.158421	14.568431	0.007975
11	-2581.711672	0.015180	0.001042
12	-2581.726852	-0.000000	-0.000000

2.8 General problems

Step 0 Given \mathbf{x}^1 , set $k := 1$, $\mathbf{d}_1 = -\nabla f(\mathbf{x}^0)$

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Step 1 If $\|\nabla f(\mathbf{x}^k)\| \leq \epsilon$, stop; else:

$$\hat{\alpha}_k = \operatorname{argmin}_{\alpha} f(\mathbf{x}^k + \alpha \mathbf{d}^k) = \frac{-\nabla f(\mathbf{x}^k)' \mathbf{d}_k}{\mathbf{d}_k' Q \mathbf{d}_k}$$

$$\mathbf{x}^{k+1} = \mathbf{x}^k + \hat{\alpha}_k \mathbf{d}_k$$

$$\mathbf{d}_{k+1} = -\nabla f(\mathbf{x}^{k+1}) + \lambda_k \mathbf{d}_k$$

$$\lambda_k = \frac{\|\nabla f(\mathbf{x}^{k+1})\|}{\|\nabla f(\mathbf{x}^k)\|}$$

Step 2 $k \leftarrow k + 1$, goto Step 1

3 Necessary Optimality Conditions

3.1 Nonlinear Optimization

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$$\begin{aligned} & \min f(\mathbf{x}) \\ \text{s.t. } & g_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, p \\ & h_i(\mathbf{x}) = 0, \quad i = 1, \dots, m \end{aligned}$$

$$P = \{\mathbf{x} \mid g_j(\mathbf{x}) \leq 0, \quad j = 1, \dots, p, \\ h_i(\mathbf{x}) = 0, \quad i = 1, \dots, m\}$$

3.2 The KKT conditions

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Discovered by Karush-Kuhn-Tucker in 1950's.

Theorem

If

- $\bar{\mathbf{x}}$ is local minimum of P
- $I = \{j \mid g_j(\bar{\mathbf{x}}) = 0\}$, set of tight constraints
- Constraint qualification condition (CQC): The vectors $\nabla g_j(\bar{\mathbf{x}})$, $j \in I$ and $\nabla h_i(\bar{\mathbf{x}})$, $i = 1, \dots, m$, are linearly independent

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Then, there exist vectors (\mathbf{u}, \mathbf{v}) :

1. $\nabla f(\bar{\mathbf{x}}) + \sum_{j=1}^p u_j \nabla g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m v_i \nabla h_i(\bar{\mathbf{x}}) = \mathbf{0}$
2. $\mathbf{u} \geq \mathbf{0}$
3. $u_j g_j(\bar{\mathbf{x}}) = 0, \quad j = 1, \dots, p$

3.3 Some Intuition from LO

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Linearize the functions in the neighborhood of the solution $\bar{\mathbf{x}}$. Problem becomes:

$$\begin{aligned} \min \quad & f(\bar{\mathbf{x}}) + \nabla f(\bar{\mathbf{x}})'(\mathbf{x} - \bar{\mathbf{x}}) \\ \text{s.t.} \quad & g_j(\bar{\mathbf{x}}) + \nabla g_j(\bar{\mathbf{x}})'(\mathbf{x} - \bar{\mathbf{x}}) \leq \mathbf{0}, \quad j \in I \\ & h_i(\bar{\mathbf{x}}) + \nabla h_i(\bar{\mathbf{x}})'(\mathbf{x} - \bar{\mathbf{x}}) = \mathbf{0}, \quad i = 1, \dots, m \end{aligned}$$

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This is a LO problem. Dual feasibility:

$$\sum_{j \in I} \hat{u}_j \nabla g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m \hat{v}_i \nabla h_i(\bar{\mathbf{x}}) = \nabla f(\bar{\mathbf{x}}), \quad \hat{u}_j \leq 0$$

Change to $u_j = -\hat{u}_j$, $v_i = -\hat{v}_i$ to obtain:

$$\nabla f(\bar{\mathbf{x}}) + \sum_{j=1}^p u_j \nabla g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m v_i \nabla h_i(\bar{\mathbf{x}}) = \mathbf{0}, \quad u_j \geq 0$$

3.4 Example 1

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$$\begin{aligned} \min \quad & f(\mathbf{x}) = (x_1 - 12)^2 + (x_2 + 6)^2 \\ \text{s.t.} \quad & h_1(\mathbf{x}) = 8x_1 + 4x_2 - 20 = 0 \\ & g_1(\mathbf{x}) = x_1^2 + x_2^2 + 3x_1 - 4.5x_2 - 6.5 \leq 0 \\ & g_2(\mathbf{x}) = (x_1 - 9)^2 + x_2^2 - 64 \leq 0 \end{aligned}$$

$$\bar{\mathbf{x}} = (2, 1)'; \quad g_1(\bar{\mathbf{x}}) = 0, \quad g_2(\bar{\mathbf{x}}) = -14, \quad h_1(\bar{\mathbf{x}}) = 0.$$

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- $I = \{1\}$
- $\nabla f(\bar{\mathbf{x}}) = (-20, 14)'; \quad \nabla g_1(\bar{\mathbf{x}}) = (7, -2.5)'$
- $\nabla g_2(\bar{\mathbf{x}}) = (-14, 2)'; \quad \nabla h_1(\bar{\mathbf{x}}) = (8, 4)'$
- $u_1 = 4, u_2 = 0, v_1 = -1$
- $\nabla g_1(\bar{\mathbf{x}}), \nabla h_1(\bar{\mathbf{x}})$ linearly independent
- $\nabla f(\bar{\mathbf{x}}) + u_1 \nabla g_1(\bar{\mathbf{x}}) + u_2 \nabla g_2(\bar{\mathbf{x}}) + v_1 \nabla h_1(\bar{\mathbf{x}}) = \mathbf{0}$

$$\begin{pmatrix} -20 \\ 14 \end{pmatrix} + 4 \begin{pmatrix} 7 \\ -2.5 \end{pmatrix} + 0 \begin{pmatrix} -14 \\ 2 \end{pmatrix} + (-1) \begin{pmatrix} 8 \\ 4 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

3.5 Example 2

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$$\begin{aligned} \max \quad & \mathbf{x}' Q \mathbf{x} \\ \text{s.t.} \quad & \mathbf{x}' \mathbf{x} \leq 1 \end{aligned}$$

Q arbitrary; Not a convex optimization problem.

$$\begin{aligned} \min \quad & -\mathbf{x}' Q \mathbf{x} \\ \text{s.t.} \quad & \mathbf{x}' \mathbf{x} \leq 1 \end{aligned}$$

3.5.1 KKT

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$$\begin{aligned} -2Q\mathbf{x} + 2u\mathbf{x} &= \mathbf{0} \\ \mathbf{x}' \mathbf{x} &\leq 1 \\ u &\geq 0 \\ u(1 - \mathbf{x}' \mathbf{x}) &= 0 \end{aligned}$$

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3.5.2 Solutions of KKT

- $\bar{\mathbf{x}} = \mathbf{0}, \bar{u} = 0$, Obj = 0.
- $\bar{\mathbf{x}} \neq \mathbf{0} \Rightarrow Q\bar{\mathbf{x}} = \bar{u}\bar{\mathbf{x}} \Rightarrow \bar{\mathbf{x}}$ eigenvector of Q with non-negative eigenvalue \bar{u} .
- $\bar{\mathbf{x}}' Q \bar{\mathbf{x}} = \bar{u} \bar{\mathbf{x}}' \bar{\mathbf{x}} = \bar{u}$.
- Thus, pick the largest nonnegative eigenvalue \hat{u} of Q . The solution is the corresponding eigenvector $\hat{\mathbf{x}}$ normalized such that $\hat{\mathbf{x}}' \hat{\mathbf{x}} = 1$. If all eigenvalues are negative, $\hat{\mathbf{x}} = \mathbf{0}$.

3.6 Are CQC Necessary?

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$$\begin{aligned} \min \quad & x_1 \\ \text{s.t.} \quad & x_1^2 - x_2 \leq 0 \\ & x_2 = 0 \end{aligned}$$

Feasible space is $(0, 0)$.

KKT:

$$\begin{pmatrix} 1 \\ 0 \end{pmatrix} + u \begin{pmatrix} 2x_1 \\ -1 \end{pmatrix} + v \begin{pmatrix} 0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \end{pmatrix}$$

KKT multipliers do not exist, while still $(0, 0)'$ is local minimum. Check $\nabla g_1(0, 0)$ and $\nabla h_1(0, 0)$.

3.7 Constrained Qualification

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Slater condition: There exists an \mathbf{x}^0 such that $g_j(\mathbf{x}^0) < 0$, $j = 1, \dots, p$, and $h_i(\mathbf{x}^0) = 0$ for all $i = 1, \dots, m$.

Theorem Under the Slater condition the KKT conditions are necessary.

4 Sufficient Optimality Conditions

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Theorem If

- $\bar{\mathbf{x}}$ feasible for P
- Feasible set is P is convex and $f(\mathbf{x})$ convex
- There exist vectors (\mathbf{u}, \mathbf{v}) , $\mathbf{u} \geq \mathbf{0}$:

$$\nabla f(\bar{\mathbf{x}}) + \sum_{j=1}^p u_j \nabla g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m v_i \nabla h_i(\bar{\mathbf{x}}) = \mathbf{0}$$

$$u_j g_j(\bar{\mathbf{x}}) = 0, \quad j = 1, \dots, p$$

Then, $\bar{\mathbf{x}}$ is a global minimum of P .

4.1 Proof

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- Let $\mathbf{x} \in P$. Then $(1 - \lambda)\bar{\mathbf{x}} + \lambda\mathbf{x} \in P$ for $\lambda \in [0, 1]$.
- $g_j((\bar{\mathbf{x}} + \lambda(\mathbf{x} - \bar{\mathbf{x}})) \leq 0 \Rightarrow$

$$\nabla g_j(\bar{\mathbf{x}})'(\mathbf{x} - \bar{\mathbf{x}}) \leq 0$$

- Similarly, $h_i(\bar{\mathbf{x}} + \lambda(\mathbf{x} - \bar{\mathbf{x}})) \leq 0 \Rightarrow$

$$\nabla h_i(\bar{\mathbf{x}})'(\mathbf{x} - \bar{\mathbf{x}}) = 0$$

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- Thus,

$$\begin{aligned} \nabla f(\bar{\mathbf{x}})'(\mathbf{x} - \bar{\mathbf{x}}) &= \\ -\left(\sum_{j=1}^p u_j \nabla g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m v_i \nabla h_i(\bar{\mathbf{x}})\right)'(\mathbf{x} - \bar{\mathbf{x}}) &\geq 0 \\ \Rightarrow f(\mathbf{x}) &\geq f(\bar{\mathbf{x}}). \end{aligned}$$

5 Convex Optimization

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- The KKT conditions are always necessary under CQC.
- The KKT conditions are sufficient for convex optimization problems.
- The KKT conditions are necessary and sufficient for convex optimization problems under CQC.
- $\min f(\mathbf{x})$ s.t. $\mathbf{A}\mathbf{x} = \mathbf{b}$, $\mathbf{x} \geq \mathbf{0}$, $f(\mathbf{x})$ convex, KKT are necessary and sufficient even without CQC.

5.0.1 Separating hyperplanes

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Theorem Let S be a nonempty closed convex subset of \Re^n and let $\mathbf{x}^* \in \Re^n$ be a vector that does not belong to S . Then, there exists some vector $\mathbf{c} \in \Re^n$ such that $\mathbf{c}'\mathbf{x}^* < \mathbf{c}'\mathbf{x}$ for all $\mathbf{x} \in S$.

Proof in BT, p.170

5.1 Sketch of the Proof under convexity

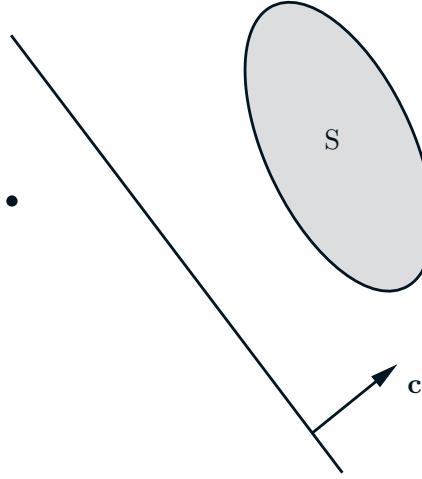
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- Suppose $\bar{\mathbf{x}}$ is a local (and thus global) optimal solution.
- $f(\mathbf{x}) < f(\bar{\mathbf{x}})$, $g_j(\mathbf{x}) \leq \mathbf{0}$, $j = 1, \dots, p$, $h_i(\mathbf{x}) = \mathbf{0}$, $i = 1, \dots, m$ is infeasible.
- Let $U = \{(u_0, \mathbf{u}, \mathbf{v}) \mid \text{there exists } \mathbf{x} : f(\mathbf{x}) < u_0, g_j(\mathbf{x}) \leq u_j, h_i(\mathbf{x}) = v_i\}$.
- $(f(\bar{\mathbf{x}}), \mathbf{0}, \mathbf{0}) \notin U$.
- U convex.

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- By separating hyperplane theorem, there is a vector $(c_0, \mathbf{c}, \mathbf{d})$:

$$c_0 u_0 + \sum_{j=1}^p c_j u_j + \sum_{i=1}^m d_i v_i > c_0 f(\bar{\mathbf{x}}) \quad \forall (u_0, \mathbf{u}, \mathbf{v}) \in U.$$



- $c_0 \geq 0$ and $c_j \geq 0$ for $j \in I$ (constraint $g_j(\bar{\mathbf{x}}) \leq 0$ tight). Why?
If $(u_0, \mathbf{u}, \mathbf{v}) \in U$, then $(u_0 + \lambda, \mathbf{u}, \mathbf{v}) \in U$ for $\lambda \geq 0$. Thus,

$$\forall \lambda \geq 0, \lambda c_0 + \sum_{j=1}^p c_j u_j + \sum_{i=1}^m d_i v_i > c_0 f(\bar{\mathbf{x}}) \Rightarrow c_0 \geq 0.$$

- Select $(u_0, \mathbf{u}, \mathbf{v}) = (f(\mathbf{x}) + \lambda, g_1(\mathbf{x}), \dots, g_p(\mathbf{x}), h_1(\mathbf{x}), \dots, h_m(\mathbf{x})) \in U$
- $c_0(f(\mathbf{x}) + \lambda) + \sum_{j=1}^p c_j g_j(\mathbf{x}) + \sum_{i=1}^m d_i h_i(\mathbf{x}) > c_0 f(\bar{\mathbf{x}})$
- Take $\lambda \rightarrow 0$:

$$c_0 f(\mathbf{x}) + \sum_{j=1}^p c_j g_j(\mathbf{x}) + \sum_{i=1}^m d_i h_i(\mathbf{x}) \geq c_0 f(\bar{\mathbf{x}})$$

- $c_0 > 0$ (constrained qualification needed here).
- $f(\mathbf{x}) + \sum_{j=1}^p u_j g_j(\mathbf{x}) + \sum_{i=1}^m v_i h_i(\mathbf{x}) \geq f(\bar{\mathbf{x}}), \quad u_j \geq 0$
- $f(\bar{\mathbf{x}}) + \sum_{j=1}^p u_j g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m v_i h_i(\bar{\mathbf{x}}) \leq f(\bar{\mathbf{x}}) \leq f(\mathbf{x}) + \sum_{j=1}^p u_j g_j(\mathbf{x}) + \sum_{i=1}^m v_i h_i(\mathbf{x}).$

- Thus,

$$f(\bar{\mathbf{x}}) = \min \left(f(\mathbf{x}) + \sum_{j=1}^p u_j g_j(\mathbf{x}) + \sum_{i=1}^m v_i h_i(\mathbf{x}) \right)$$

$$\sum_{j=1}^p u_j g_j(\bar{\mathbf{x}}) = 0 \Rightarrow u_j g_j(\bar{\mathbf{x}}) = 0$$

- Unconstrained optimality conditions:

$$\nabla f(\bar{\mathbf{x}}) + \sum_{j=1}^p u_j \nabla g_j(\bar{\mathbf{x}}) + \sum_{i=1}^m v_i \nabla h_i(\bar{\mathbf{x}}) = \mathbf{0}$$

6 Applications

6.1 Linear Optimization

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$$\begin{aligned} & \min \quad \mathbf{c}' \mathbf{x} \\ \text{s.t.} \quad & \mathbf{A} \mathbf{x} = \mathbf{b} \\ & \mathbf{x} \geq \mathbf{0} \end{aligned}$$

$$\begin{aligned} & \min \quad \mathbf{c}' \mathbf{x} \\ \text{s.t.} \quad & \mathbf{A} \mathbf{x} - \mathbf{b} = \mathbf{0} \\ & -\mathbf{x} \leq \mathbf{0} \end{aligned}$$

6.1.1 KKT

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$$\begin{aligned} \mathbf{c} + \mathbf{A}' \hat{\mathbf{u}} - \mathbf{v} &= \mathbf{0} \\ \mathbf{v} &\geq \mathbf{0} \\ v_j x_j &= 0 \\ \mathbf{A} \mathbf{x} - \mathbf{b} &= \mathbf{0} \\ \mathbf{x} &\geq \mathbf{0} \end{aligned}$$

$$\mathbf{u} = -\hat{\mathbf{u}}$$

$$\begin{aligned} \mathbf{A}' \mathbf{u} &\leq \mathbf{c} && \text{dual feasibility} \\ (c_j - \mathbf{A}'_j \mathbf{u}) x_j &= 0 && \text{complementarity} \\ \mathbf{A} \mathbf{x} - \mathbf{b} &= \mathbf{0} && \text{primal feasibility} \\ \mathbf{x} &\geq \mathbf{0} && \text{primal feasibility} \end{aligned}$$

6.2 Portfolio Optimization

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\mathbf{x} =weights of the portfolio

$$\begin{aligned} \max \quad & \mathbf{r}'\mathbf{x} - \lambda \frac{1}{2}\mathbf{x}'\mathbf{Q}\mathbf{x} \\ \text{s.t.} \quad & \mathbf{e}'\mathbf{x} = 1 \end{aligned}$$

$$\begin{aligned} \min \quad & -\mathbf{r}'\mathbf{x} + \lambda \frac{1}{2}\mathbf{x}'\mathbf{Q}\mathbf{x} \\ \text{s.t.} \quad & \mathbf{e}'\mathbf{x} = 1 \end{aligned}$$

6.2.1 KKT

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$$\begin{aligned} -\mathbf{r} + \lambda \mathbf{Q}\mathbf{x} + u\mathbf{e} &= 0 \\ \mathbf{x} &= \frac{1}{\lambda} \mathbf{Q}^{-1}(\mathbf{r} - u\mathbf{e}) \\ \mathbf{e}'\mathbf{x} = 1 \Rightarrow \mathbf{e}'\mathbf{Q}^{-1}(\mathbf{r} - u\mathbf{e}) &= \lambda \\ u &= \frac{\mathbf{e}'\mathbf{Q}^{-1}\mathbf{r} - \lambda}{\mathbf{e}'\mathbf{Q}^{-1}\mathbf{e}} \end{aligned}$$

As λ changes, tradeoff of risk and return changes. The allocation changes as well. This is the essence of modern portfolio theory.

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