1.204 Quiz 1 Spring 2010 Solutions

1. Data model (35 points)



2. Database/SQL. (30 points)

Based on the data model in question 2, you construct a database where the tables, attributes and relationships correspond exactly to the data model structure. You are now asked to write the following SQL queries against your database.

a. List the name and vendor ID of all vendors that supply at least one assembly to MITG. List each vendor only once. (10 points)

SELECT DISTINCT Vendor.VendorName, Vendor.VendorID FROM Vendor, VendorAssembly WHERE Vendor.VendorID= VendorAssembly.VendorID;

> b. List the product ID and glider description of all high performance hang gliders with a wholesale sales type. Assume the code for high performance is "H" and for wholesale is "W" (10 points)

SELECT Glider.ProductID, Glider.GliderDesc FROM Product, Glider WHERE Product.ProductID= Glider.ProductID AND Product.SalesType= "W" AND Glider.GliderType= "H";

c. List the name and vendor ID of all vendors whose average assembly price is above \$50. (10 points)

SELECT Vendor.VendorName, Vendor.VendorID FROM Vendor, VendorAssembly WHERE Vendor.VendorID = VendorAssembly.VendorID GROUP BY Vendor.VendorName, Vendor.VendorID HAVING (AVG(VendorAssembly.VendorAssemblyPrice)>50);

3. Algorithm design. (35 points)

Each job in the figure below starts at a time s_i and finishes at a time f_i . Two jobs are compatible if they do not overlap. Find the maximum subset of compatible jobs. Give an algorithm that solves this problem.



Figure by MIT OpenCourseWare.

a. What kind of algorithm is it: divide and conquer, greedy, etc.?

Greedy.

b. Write the algorithm in pseudocode or Java; comment or explain it so that it is easy to understand. Sort jobs by earliest finish time. Try to add each job in order.

```
return A
```

c. Informally show that your algorithm is correct.



Counterexamples for greedy algorithms that aren't quite right:

Optimal solution: Sort jobs by earliest finish time. Proof by contradiction. Select initial set of jobs in greedy. Assume a job with a later finish time is in the optimal set. But this is impossible, since it uses more of the scarce time resource and gains nothing

Figure by MIT OpenCourseWare.

1.204 Computer Algorithms in Systems Engineering Spring 2010

For information about citing these materials or our Terms of Use, visit: http://ocw.mit.edu/terms.